# Introduction to OpenACC

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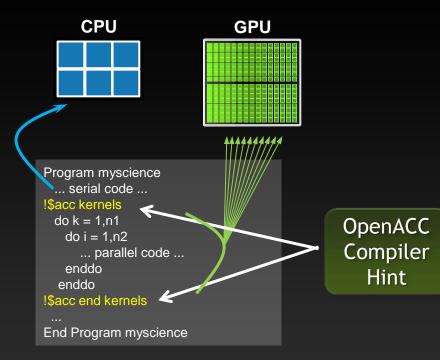
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# What is OpenACC?

It is a directive based standard to allow developers to take advantage of accelerators such as GPUs from NVIDIA and AMD, Intel's Xeon Phi, FPGAs, and even DSP chips.



# **Directives**



Simple compiler hints from coder.

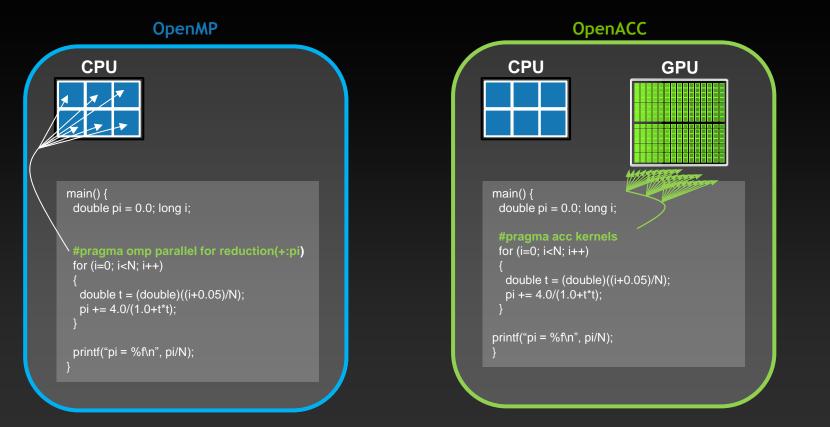
Compiler generates parallel threaded code.

Ignorant compiler just sees some comments.

Your original Fortran or C code



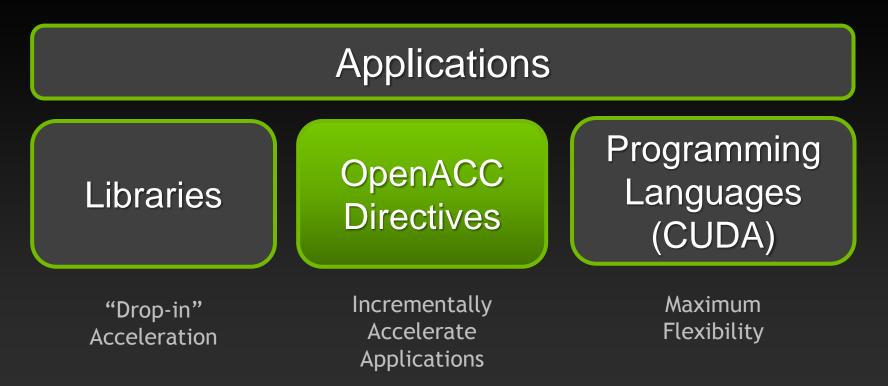
## Familiar to OpenMP Programmers



#### More on this later!



# How Else Would We Accelerate Applications?





# Key Advantages Of This Approach

- High-level. No involvement of OpenCL, CUDA, etc.
- Single source. No forking off a separate GPU code. Compile the same program for accelerators or serial, non-GPU programmers can play along.
- Efficient. Experience shows very favorable comparison to low-level implementations of same algorithms.
- Performance portable. Supports GPU accelerators and co-processors from multiple vendors, current and future versions.
- Incremental. Developers can port and tune parts of their application as resources and profiling dictates. No wholesale rewrite required. Which can be <u>quick.</u>



# **A Few Cases**

Designing circuits for quantum computing

UIST, Macedonia

Reading DNA nucleotide sequences Shanghai JiaoTong University



## 4 directives

16x faster

HydroC- Galaxy Formation

PRACE Benchmark Code, CAPS



1 week

3x faster



1 week

40x faster

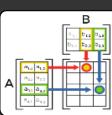
**Real-time Derivative Valuation** 

Opel Blue, Ltd

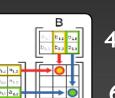


Few hours

70x faster







#### 4 directives





Extracting image features in realtime

Aselsan

Matrix Matrix Multiply

3 directives

4.1x faster

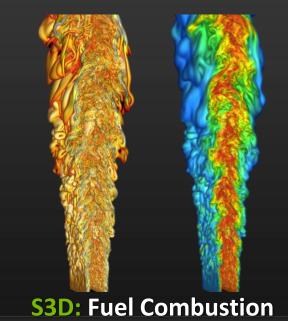
## **A Champion Case**

# 4x FasterJaguarTitan42 days10 days

## Modified <1% Lines of Code

15 PF! One of fastest simulations ever!

Design alternative fuels with up to 50% higher efficiency

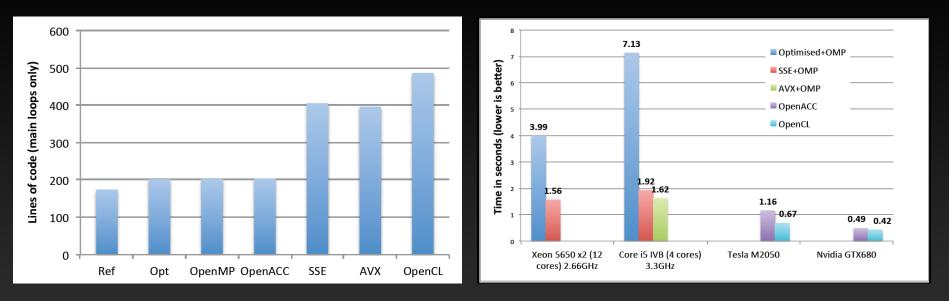




#### The Portland Group®

## **Comparison to Alternatives**

#### Lattice-Boltzmann Example





## **Broad Accelerator Support**

- Xeon Phi support already in CAPS. Demonstrated and soon to be release for PGI.
- AMD line of accelerated processing units (APUs) as well as the AMD line of discrete GPUs for preliminary PGI support.
- Carma a hybrid platform based on ARM Cortex-A9 quad core and an NVIDIA Quadro® 1000M GPU.
- NVIDIA...



## **NVIDIA Rules**

or writes the rules. They have been the foremost supporter of GPU computing for much of the past decade, and have earned the focus of this workshop. We are using NVIDIA GPUs as our platform and our touchstone because:

- They are proven
- Well understood
- Best bang for buck if you want to buy an accelerator
- Excellent support by vendor and community
- It is the basis for our leading edge platform, Keeneland
- It will not be going obsolete any time soon
- NVIDIA recently acquired PGI. That gave us a slight preference for the PGI compiler over the Cray one. Both are available on Blue Waters.



## **True Standard**

Full OpenACC 1.0 and 2.0 Specifications available online

http://www.openacc-standard.org

- Quick reference card also available
- Implementations available now from PGI, Cray, and CAPS.
- GCC version of OpenACC in 4.9x and standard in 5.0 (official release early April).

#### The OpenACC<sup>™</sup> API QUICK REFERENCE GUIDE

The OpenACC Application Program Interface describes a collection of compiler directives to specify loops and regions of code in standard C, C++ and Fortran to be offloaded from a host CPU to an attached accelerator, providing portability across operating systems, host CPUs and accelerators.

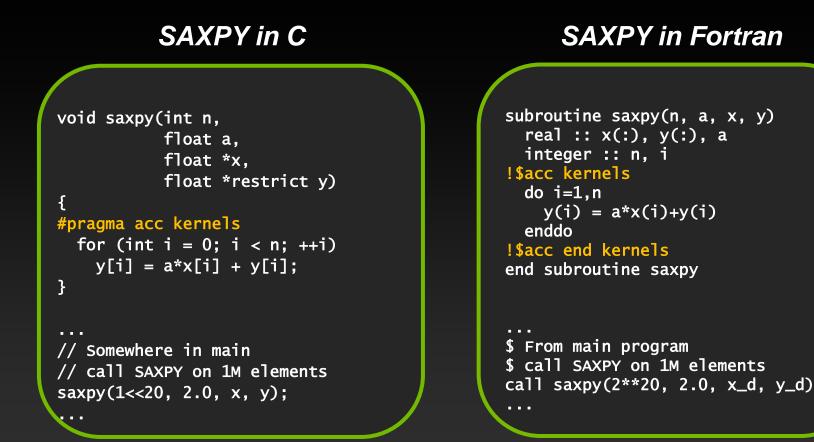
Most OpenACC directives apply to the immediately following structured block or loop; a structured block is a single statement or a compound statement (C or C++) or a sequence of statements (Cortran) with a single entry point at the top and a single exit at the bottom.



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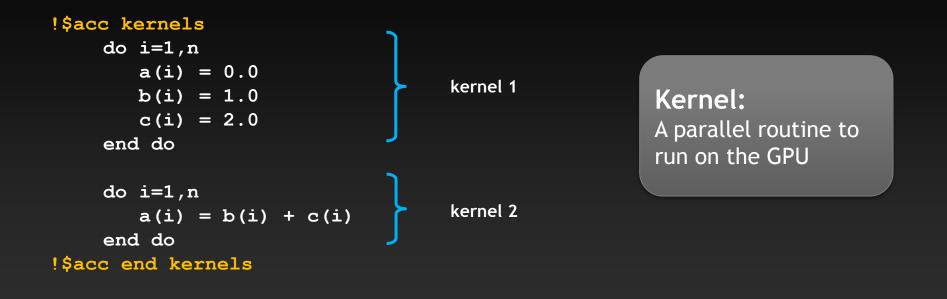
## A Simple Example: SAXPY





## kernels: Our first OpenACC Directive

We request that each loop execute as a separate *kernel* on the GPU. This is an incredibly powerful directive.





## **General Directive Syntax and Scope**

Fortran

!\$acc kernels [clause ...]
 structured block
!\$acc end kernels

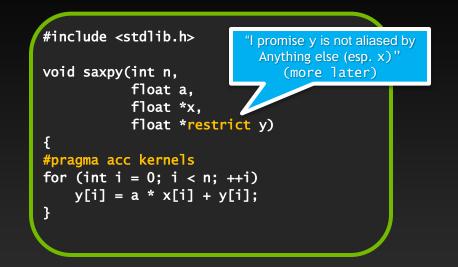
#pragma acc kernels [clause ...]
{
 structured block
}

I may indent the directives at the natural code indentation level for readability. It is a common practice to always start them in the first column (ala #define/#ifdef). Either is fine with C or Fortran 90 compilers.



## **Complete SAXPY Example Code**

```
int main(int argc, char **argv)
 int N = 1 << 20; // 1 million floats
 if (argc > 1)
   N = atoi(argv[1]);
 float *x = (float*)malloc(N * sizeof(float));
  float *y = (float*)malloc(N * sizeof(float));
  for (int i = 0; i < N; ++i) {
   x[i] = 2.0f;
   y[i] = 1.0f;
  }
 saxpy(N, 3.0f, x, y);
  return 0;
```



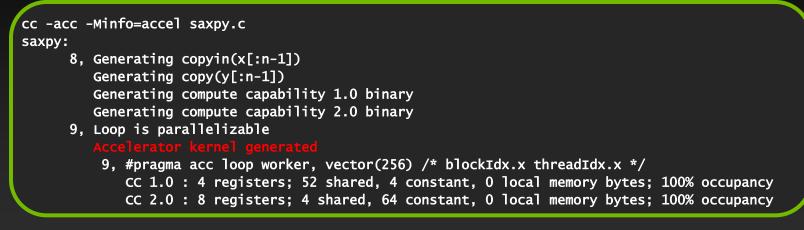


## **Compile and Run**

C: cc -acc -Minfo=accel saxpy.c

Fortran: ftn –acc –Minfo=accel saxpy.f90

#### **Compiler Output**



Run: aprun -n1 a.out



## Compare: Partial CUDA C SAXPY Code Just the subroutine

```
global void saxpy kernel( float a, float* x, float* y, int n ){
 int i;
 i = blockIdx.x*blockDim.x + threadIdx.x;
 if(i \le n) x[i] = a*x[i] + y[i];
void saxpy( float a, float* x, float* y, int n ){
  float *xd, *yd;
  cudaMalloc( (void**)&xd, n*sizeof(float) );
  cudaMalloc( (void**)&yd, n*sizeof(float) ); cudaMemcpy( xd, x, n*sizeof(float),
                     cudaMemcpyHostToDevice );
  cudaMemcpy( vd, v, n*sizeof(float),
                     cudaMemcpyHostToDevice );
  saxpy kernel<<< (n+31)/32, 32 >>>( a, xd, yd, n );
  cudaMemcpy( x, xd, n*sizeof(float),
                     cudaMemcpyDeviceToHost );
  cudaFree( xd ); cudaFree( yd );
```



## Compare: Partial CUDA Fortran SAXPY Code Just the subroutine

```
module kmod
 use cudafor
contains
 attributes(global) subroutine saxpy kernel(A,X,Y,N)
  real(4), device :: A, X(N), Y(N)
  integer, value :: N
  integer :: i
  i = (blockidx%x-1)*blockdim%x + threadidx%x
  if(i \le N) X(i) = A*X(i) + Y(i)
 end subroutine
end module
 subroutine saxpy(A, X, Y, N)
  use kmod
  real(4) :: A, X(N), Y(N)
  integer :: N
  real(4), device, allocatable, dimension(:):: &
                 Xd, Yd
  allocate(Xd(N), Yd(N))
  Xd = X(1:N)
  Yd = Y(1:N)
  call saxpy kernel \langle \langle (N+31)/32, 32 \rangle \rangle \langle A, Xd, Yd, N \rangle
  X(1:N) = Xd
  deallocate(Xd, Yd)
 end subroutine
```



## Again: Complete SAXPY Example Code

#### Main Code

```
int main(int argc, char **argv)
 int N = 1 << 20; // 1 million floats
 if (argc > 1)
   N = atoi(argv[1]);
 float *x = (float*)malloc(N * sizeof(float));
  float *y = (float*)malloc(N * sizeof(float));
  for (int i = 0; i < N; ++i) {
   x[i] = 2.0f;
   y[i] = 1.0f;
  }
 saxpy(N, 3.0f, x, y);
  return 0;
```

## **Entire Subroutine**

#include <stdlib.h>

```
void saxpy(int n,
float a,
float *x,
float *restrict y)
```

```
#pragma acc kernels
for (int i = 0; i < n; ++i)
    y[i] = a * x[i] + y[i];</pre>
```



# **Big Difference!**

- With CUDA, we changed the structure of the old code. Non-CUDA programmers can't understand new code. It is not even ANSI standard code.
- We have separate sections for the host code, and the GPU code. Different flow of code. Serial path now gone forever.
- Where did these "32's" and other mystery variables come from? This is a clue that we have some hardware details to deal with here.
- Exact same situation as assembly used to be. How much hand-assembled code is still being written in HPC now that compilers have gotten so efficient?



## This looks easy! Too easy...

- If it is this simple, why don't we just throw kernel in front of every loop?
- Better yet, why doesn't the compiler do this for me?

The answer is that there are two general issues that prevent the compiler from being able to just automatically parallelize every loop.

- Data Dependencies in Loops
  - Data Movement

The compiler needs your higher level perspective (in the form of directive hints) to get correct results, and reasonable performance.



## **Data Dependencies**

Most directive based parallelization consists of splitting up big do/for loops into independent chunks that the many processors can work on simultaneously.

Take, for example, a simple for loop like this:

for(index=0, index<1000000,index++)
 Array[index] = 4 \* Array[index];</pre>

When run on 1000 processors, it will execute something like this...



## No Data Dependency





## **Data Dependency**

But what if the loops are not entirely independent?

Take, for example, a similar loop like this:

## 

This is perfectly valid serial code.



## **Data Dependency**

Now Processor 2, in trying to calculate its first iteration...

needs the result of Processor 1's last iteration. If we want the correct ("same as serial") result, we need to wait until processor 1 finishes. Likewise for processors 3, 4, ...



## **Data Dependencies**

That is a data dependency. If the compiler even <u>suspects</u> that there is a data dependency, it will, for the sake of correctness, refuse to parallelize that loop.

11, Loop carried dependence of 'Array' prevents parallelization Loop carried backward dependence of 'Array' prevents vectorization

As large, complex loops are quite common in HPC, especially around the most important parts of your code, the compiler will often balk most when you most need a kernel to be generated. What can you do?



## **Data Dependencies**

- Rearrange your code to make it more obvious to the compiler that there is not really a data dependency.
- Eliminate a real dependency by changing your code.
  - There is a common bag of tricks developed for this as this issue goes back 40 years in HPC. Many are quite trivial to apply.
  - The compilers have gradually been learning these themselves.
- Override the compiler's judgment (independent clause) at the risk of invalid results. Misuse of restrict has similar consequences.



## C Detail: the restrict keyword

- Standard C (as of C99).
- Important for optimization of serial as well as OpenACC and OpenMP code.
- Promise given by the programmer to the compiler for a pointer float \*restrict ptr

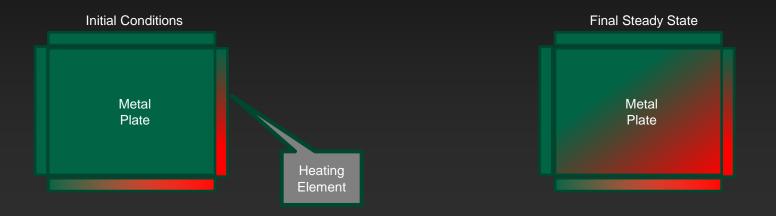
Meaning: "for the lifetime of ptr, only it or a value directly derived from it (such as ptr + 1) will be used to access the object to which it points"

- Limits the effects of pointer aliasing
- OpenACC compilers often require restrict to determine independence
  - Otherwise the compiler can't parallelize loops that access ptr
  - Note: if programmer violates the declaration, behavior is undefined



## **Our Foundation Exercise: Laplace Solver**

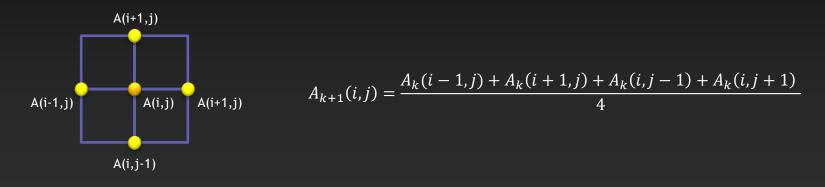
- I've been using this for MPI, OpenMP and now OpenACC. It is a great simulation problem, not rigged for OpenACC.
- In this most basic form, it solves the Laplace equation:  $abla^2 f(x,y) = 0$
- The Laplace Equation applies to many physical problems, including:
  - Electrostatics
  - Fluid Flow
  - Temperature
- For temperature, it is the Steady State Heat Equation:





## **Exercise Foundation: Jacobi Iteration**

- The Laplace equation on a grid states that each grid point is the average of it's neighbors.
- We can iteratively converge to that state by repeatedly computing new values at each point from the average of neighboring points.
- We just keep doing this until the difference from one pass to the next is small enough for us to tolerate.





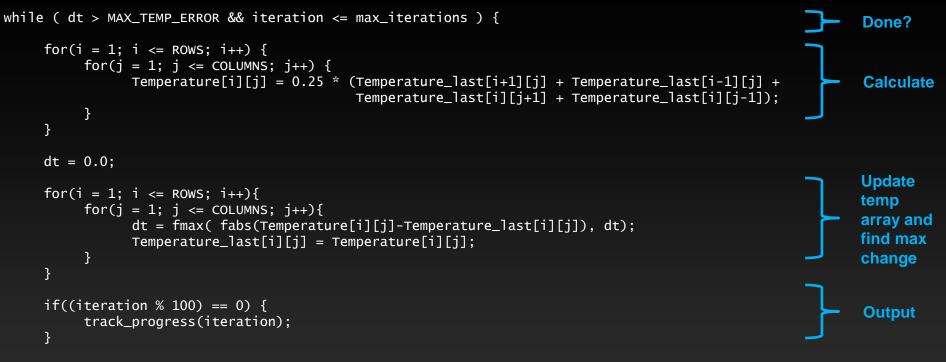
## Serial Code Implementation







## Serial C Code (kernel)



iteration++;



## Serial C Code Subroutines

#### void initialize(){

```
int i,j;
for(i = 0; i <= ROWS+1; i++){</pre>
    for (j = 0; j \le COLUMNS+1; j++)
        Temperature_last[i][j] = 0.0;
}
// these boundary conditions never change throughout run
// set left side to 0 and right to a linear increase
for(i = 0; i <= ROWS+1; i++) {</pre>
    Temperature_last[i][0] = 0.0;
    Temperature_last[i][COLUMNS+1] = (100.0/ROWS)*i;
}
// set top to 0 and bottom to linear increase
for(j = 0; j <= COLUMNS+1; j++) {</pre>
    Temperature_last[0][j] = 0.0;
    Temperature_last[ROWS+1][j] = (100.0/COLUMNS)*j;
```

BCs could run from 0 to ROWS+1 or from 1 to ROWS. We chose the former.

```
void track_progress(int iteration) {
    int i;
    printf("-- Iteration: %d --\n", iteration);
    for(i = ROWS-5; i <= ROWS; i++) {
        printf("[%d,%d]: %5.2f ", i, i,Temperature[i][i]);
    }
    printf("\n");
}</pre>
```



#### Whole C Code

#include <stdlib.h>
#include <stdio.h>
#include <math.h>
#include <sys/time.h>

// size of plate #define COLUMNS 1000 #define ROWS 1000

// largest permitted change in temp (This value takes about 3400 steps)
#define MAX\_TEMP\_ERROR 0.01

double Temperature[ROWS+2][COLUMNS+2]; // temperature grid double Temperature\_last[ROWS+2][COLUMNS+2]; // temperature grid from last iteration

// helper routines
void initialize();
void track\_progress(int iter);

int main(int argc, char \*argv[]) {

int i, j; // grid indexes
int max\_iterations; // number of iterations
int iteration=1; // current iteration
double dt=100; // largest change in t
struct timeval start\_time, stop\_time, elapsed\_time; // timers

printf("Maximum iterations [100-4000]?\n"); scanf("%d", &max\_iterations);

gettimeofday(&start\_time,NULL); // Unix timer

initialize();

// initialize Temp\_last including boundary conditions

```
// do until error is minimal or until max steps
while ( dt > MAX_TEMP_ERROR && iteration <= max_iterations ) {</pre>
```

```
dt = 0.0; // reset largest temperature change
```

```
// copy grid to old grid for next iteration and find latest dt
for(i = 1; i <= ROWS: i++){
    for(j = 1; j <= COLUMNS; j++){
        dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
        Temperature_last[i][j] = Temperature[i][j];
    }
}
```

```
// periodically print test values
if((iteration % 100) == 0) {
    track_progress(iteration);
}
```

iteration++;

gettimeofday(&stop\_time,NULL); timersub(&stop\_time, &start\_time, &elapsed\_time); // Unix time subtract routine

printf("\nMax error at iteration %d was %f\n", iteration-1, dt); printf("Total time was %f seconds.\n", elapsed\_time.tv\_sec+elapsed\_time.tv\_usec/1000000.0);

```
// initialize plate and boundary conditions
// Temp_last is used to to start first iteration
void initialize(){
```

```
int i,j;
for(i = 0; i <= ROWS+1; i++){
    for (j = 0; j <= COLUMNS+1; j++){
        Temperature_last[i][j] = 0.0;
    }
}
```

// these boundary conditions never change throughout run

```
// set left side to 0 and right to a linear increase
for(i = 0; i <= ROWS+1; i++) {
    Temperature_last[i][0] = 0.0;
    Temperature_last[i][COLUMNS+1] = (100.0/ROWS)*i;
}
```

```
// set top to 0 and bottom to linear increase
for(j = 0; j <= COLUMNS+1; j++) {
   Temperature_last[0][j] = 0.0;
   Temperature_last[ROWS+1][j] = (100.0/COLUMNS)*j;
```

```
د
```

3

3

// print diagonal in bottom right corner where most action is
void track\_progress(int iteration) {

int i;

```
printf("------ Iteration number: %d ------\n", iteration);
for(i = ROWS-5; i <= ROWS; i++) {
    printf("[%d,%d]: %5.2f ", i, i, Temperature[i][i]);
  }
printf("\n");
```



## Serial Fortran Code (kernel)



#### enddo



#### Serial Fortran Code Subroutines

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,j

double precision, dimension(0:rows+1,0:columns+1) :: temperature\_last

 $temperature_last = 0.0$ 

!these boundary conditions never change throughout run

!set left side to 0 and right to linear increase do i=0,rows+1 temperature\_last(i,0) = 0.0 temperature\_last(i,columns+1) = (100.0/rows) \* i enddo

```
!set top to 0 and bottom to linear increase
do j=0,columns+1
    temperature_last(0,j) = 0.0
    temperature_last(rows+1,j) = ((100.0)/columns) * j
enddo
```

end subroutine initialize

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,iteration

double precision, dimension(0:rows+1,0:columns+1) :: temperature



#### Whole Fortran Code

#### program serial implicit none

!Size of plate integer, parameter :: columns=1000 integer, parameter :: rows=1000 double precision, parameter :: max\_temp\_error=0.01

integer double precision real :: i, j, max\_iterations, iteration=1
:: dt=100.0
:: start\_time, stop\_time

double precision, dimension(0:rows+1,0:columns+1) :: temperature, temperature\_last

print\*, 'Maximum iterations [100-4000]?'
read\*, max\_iterations

call cpu\_time(start\_time) !Fortran timer

call initialize(temperature\_last)

!do until error is minimal or until maximum steps
do while ( dt > max\_temp\_error .and. iteration <= max\_iterations)</pre>

#### do j=1,columns

#### enddo

dt=0.0

!copy grid to old grid for next iteration and find max change do j=1,columns do i=1,rows dt = max( abs(temperature(i,j) - temperature\_last(i,j)), dt ) temperature\_last(i,j) = temperature(i,j) enddo enddo

!periodically print test values if( mod(iteration,100).eq.0 ) then call track\_progress(temperature, iteration) endif

#### iteration = iteration+1

#### enddo

call cpu\_time(stop\_time)

print\*, 'Max error at iteration ', iteration-1, ' was ',dt
print\*, 'Total time was ',stop\_time-start\_time, ' seconds.'

end program serial

! initialize plate and boundery conditions
! temp\_last is used to to start first iteration
subroutine initialize( temperature\_last )
 implicit none

integer, parameter integer, parameter integer :: columns=1000 :: rows=1000 :: i,j

double precision, dimension(0:rows+1,0:columns+1) :: temperature\_last

 $temperature_last = 0.0$ 

!these boundary conditions never change throughout run

!set left side to 0 and right to linear increase do i=0,rows+1 temperature\_last(i,0) = 0.0 temperature\_last(i,columns+1) = (100.0/rows) \* i enddo

!set top to 0 and bottom to linear increase do j=0,columns+1 temperature\_last(0,j) = 0.0 temperature\_last(rows+1,j) = ((100.0)/columns) \* j enddo

end subroutine initialize

> integer, parameter integer, parameter integer

:: columns=1000 :: rows=1000 :: i,iteration

double precision, dimension(0:rows+1,0:columns+1) :: temperature



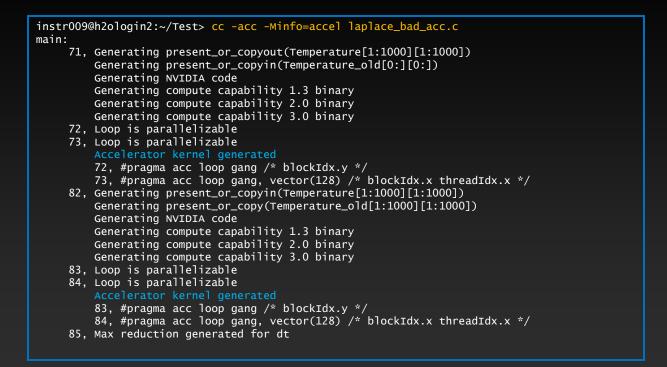
### **Exercises: General Instructions for Compiling**

- Exercises are in the "Exercises/Laplace" directory in your home directory
- Solutions are in the "Laplace/Solutions" subdirectory
- To compile cc -acc laplace.c ftn -acc laplace.f90
- This will generate the executable a.out



## Exercises: Very useful compiler option

Adding -Minfo=accel to your compile command will give you some very useful information about how well the compiler was able to honor your OpenACC directives.





#### **Exercises: General Instructions for Running**

Make sure you are in an interactive session - with idev - if you aren't already. The command prompt is your clue.

To run, use aprun:

instr003@nid25357:~> aprun -n1 a.out

You can compare against the serial code you are starting with to see what performance gains you achieve. You can compile the serial version without any extra flags (just cc or ftn), but run it as per the above. Rename your a.out's to avoid confusion.



# Exercise 1: Using kernels to parallelize the main loops (About 45 minutes)

Q: Can you get a speedup with just the kernels directives?

- 1. Edit *laplace\_serial.c/*f90
  - 1. Maybe copy your intended OpenACC version to *laplace\_acc.c* to start
  - 2. Add directives where it helps
- 2. Compile with OpenACC parallelization
  - 1. cc -acc -Minfo=accel laplace\_acc.c or ftn -acc -Minfo=accel laplace\_acc.f90
  - 2. Look at your compiler output to make sure you are having an effect
- 3. Run
  - 1. aprun -n1 a.out (Try 4000 iterations if you want a solution that converges to current tolerance)
  - 2. Serial version for baseline time
  - 3. Your OpenACC version for performance difference



### **Exercise 1 C Solution**

while ( dt > MAX\_TEMP\_ERROR && iteration <= max\_iterations ) {</pre>

}

```
Generate a GPU kernel
#pragma acc kernels
for(i = 1; i <= ROWS; i++) {</pre>
    for(j = 1; j <= COLUMNS; j++) {</pre>
        Temperature[i][j] = 0.25 * (Temperature_last[i+1][j] + Temperature_last[i-1][j] +
                                     Temperature_last[i][j+1] + Temperature_last[i][j-1]);
    }
}
dt = 0.0; // reset largest temperature change
                                                                                      Generate a GPU kernel
#pragma acc kernels
for(i = 1; i <= ROWS; i++){</pre>
    for(j = 1; j <= COLUMNS; j++){</pre>
        dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
        Temperature_last[i][i] = Temperature[i][i];
    }
}
if((iteration % 100) == 0) {
    track_progress(iteration);
}
iteration++;
```

```
PITTSBURGH
SUPERCOMPUTING
CENTER
```

### **Exercise 1 Fortran Solution**

do while ( dt > max\_temp\_error .and. iteration <= max\_iterations)</pre>

enddo

```
Generate a GPU kernel
!$acc kernels
do j=1,columns
   do i=1, rows
      temperature(i,j)=0.25*(temperature_last(i+1,j)+temperature_last(i-1,j)+ &
                             temperature_last(i,j+1)+temperature_last(i,j-1) )
   enddo
enddo
!$acc end kernels
dt=0.0
                                                                                    Generate a GPU kernel
!$acc kernels
do j=1,columns
   do i=1,rows
      dt = max( abs(temperature(i,j) - temperature_last(i,j)), dt )
      temperature_last(i,j) = temperature(i,j)
   enddo
enddo
!$acc end kernels
if( mod(iteration, 100).eq.0 ) then
   call track_progress(temperature, iteration)
endif
iteration = iteration+1
```



## Exercise 1: Compiler output (C)

instr009@h2ologin2:~/Update> cc -acc -Minfo=accel laplace\_bad\_acc.c main: 62, Generating present\_or\_copyout(Temperature[1:1000][1:1000]) Generating present\_or\_copyin(Temperature\_last[0:][0:]) Generating NVIDIA code Compiler was able to Generating compute capability 1.3 binary parallelize Generating compute capability 2.0 binary Generating compute capability 3.0 binary 63, Loop is parallelizable 64, Loop is parallelizable Accelerator kernel generated 63, #pragma acc loop gang /\* blockIdx.y \*/ 64, #pragma acc loop gang, vector(128) /\* blockIdx.x threadIdx.x \*/ 73, Generating present\_or\_copyin(Temperature[1:1000][1:1000]) Generating present\_or\_copy(Temperature\_last[1:1000][1:1000]) Generating NVIDIA code Compiler was able to Generating compute capability 1.3 binary parallelize Generating compute capability 2.0 binary Generating compute capability 3.0 binary 74. Loop is parallelizable 75, Loop is parallelizable Accelerator kernel generated 74, #pragma acc loop gang /\* blockIdx.y \*/ 75, #pragma acc loop gang, vector(128) /\* blockIdx.x threadIdx.x \*/ 76. Max reduction generated for dt



## Exercise 1: Performance

3372 steps to convergence

Execution	Time (s)	Speedup
CPU Serial (C)	36	
CPU 2 OpenMP threads	24	1.5x
CPU 4 OpenMP threads	15	2.4x
CPU 8 OpenMP threads	9.8	3.7x
CPU 16 OpenMP threads	5.0	7.2
OpenACC GPU	64	<b>0.6x</b> (0.08 vs. 16 CPU)

CPU: AMD 6276 Interlagos 8 Cores @ 2.3+GHz GPU: NVIDIA GK110 Kepler



## What's with the OpenMP?

We can compare our GPU results to the <u>best</u> the multi-core XEON CPUs can do.

If you are familiar with OpenMP, or even if you are not, you can compile and run the OpenMP enabled versions in your OpenMP directory as:

cc -mp=nonuma laplace\_omp.c or ftn -mp=nonuma laplace\_omp.f90

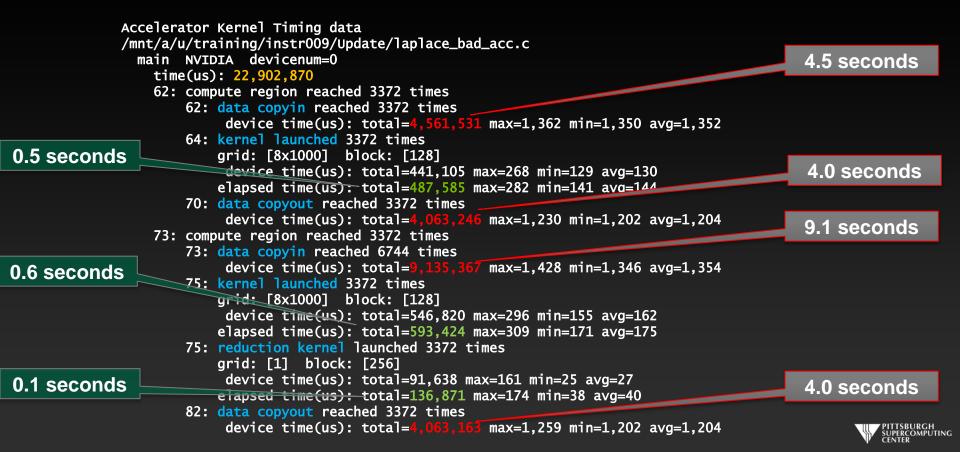
then to run on 8 threads do:

export OMP\_NUM\_THREADS=8 aprun -n 1 -d 8 a.out



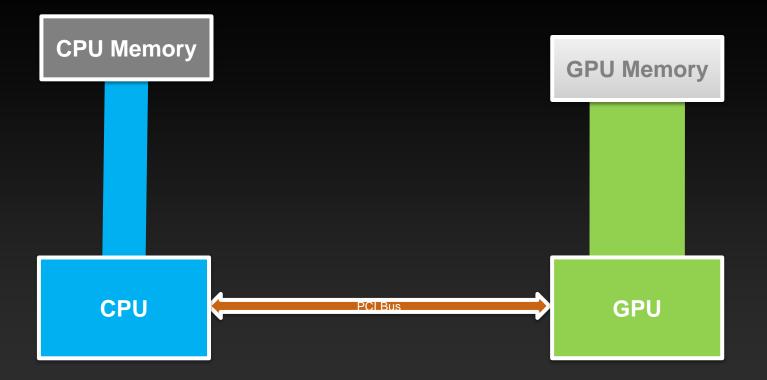
## What went wrong?

#### **export PGI\_ACC\_TIME=1** to activate profiling and run again:



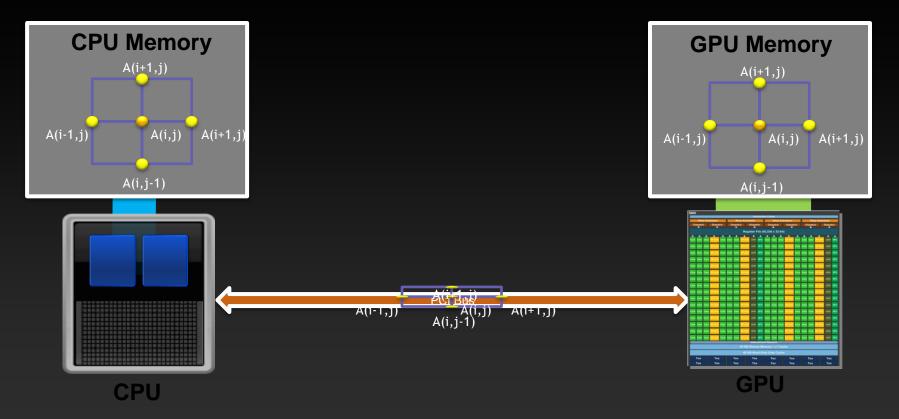
# Basic Concept

#### Simplified, but sadly true





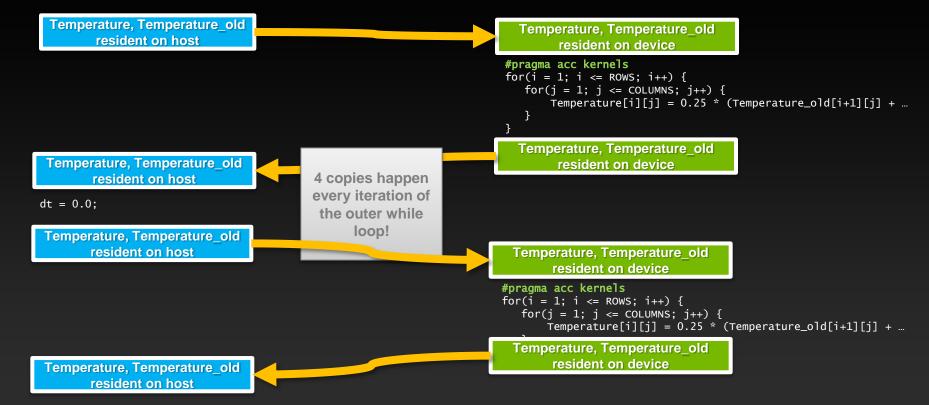
#### **Multiple Times Each Iteration**





### **Excessive Data Transfers**

#### while ( dt > MAX\_TEMP\_ERROR && iteration <= max\_iterations ) {</pre>



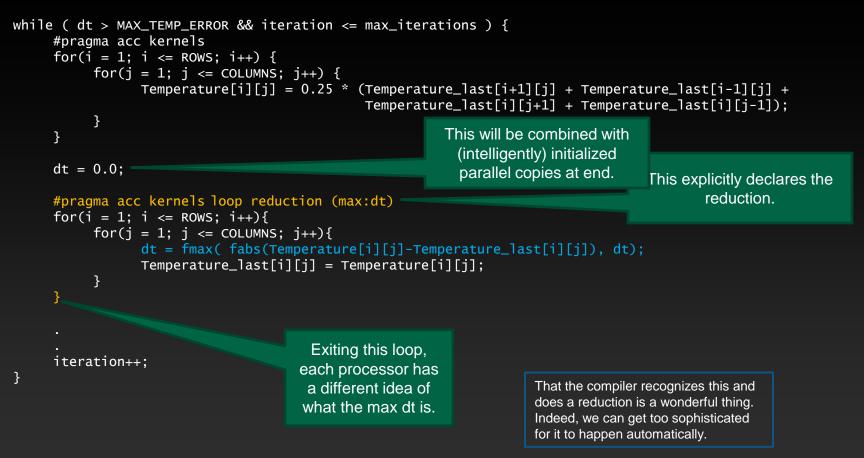


#### Data Management

The First, Most Important, and possibly Only OpenACC Optimization



#### First, about that "reduction"





## Data Construct Syntax and Scope

#### Fortran

!\$acc data [clause ...]
 structured block
!\$acc end data

#### С

#pragma acc data [clause ...]
{
 structured block
}



#### **Data Clauses**

copy( list )

Allocates memory on GPU and copies data from host to GPU when entering region and copies data to the host when exiting region. Principal use: For many important data structures in your code, this is a logical default to input, modify and return the data.

copyin(list) Allocates memory on GPU and copies data from host to GPU when
entering region.
Principal use: Think of this like an array that you would use as just
an input to a subroutine.

copyout( list ) Allocates memory on GPU and copies data to the host when exiting
region.
Principal use: A result that isn't overwriting the input data structure.

create( list ) Allocates memory on GPU but does not copy.
Principal use: Temporary arrays.



#### **Present Data Clauses**

The "present" data clauses are used when the data is already present because of a containing data region.

present( <i>list</i> )	Data is already present on GPU from another containing data region.	
	Principal use: You are calling this routine from inside a routine that	
	already has a data clause.	

present\_or\_copy
present\_or\_copyin
present\_or\_copyout
present\_or\_create

# You can't be positive that the data is present from a surrounding data region.

Principal use: A subroutine that may or may not be called from within a data region or for multi-threaded codes so that only one thread migrates data.



## Array Shaping

Compilers sometimes cannot determine the size of arrays, so we must specify explicitly using data clauses with an array "shape". The compiler will let you know if you need to do this. Sometimes, you will want to for your own efficiency reasons.

#pragma acc data copyin(a[0:size]), copyout(b[s/4:3\*s/4])

Fortran

!\$acc data copyin(a(1:size)), copyout(b(s/4:3\*s/4))

- Fortran uses start:end and C uses start:length
- Data clauses can be used on data, kernels or parallel



#### Compiler will (increasingly) often make a good guess...

```
int main(int argc, char *argv[]) {
```

```
int i;
double A[2000], B[1000], C[1000];
#pragma acc kernels
for (i=0; i<1000; i++){</pre>
```

```
A[i] = 4 * i;
B[i] = B[i] + 2;
C[i] = A[i] + 2 * B[i];
```



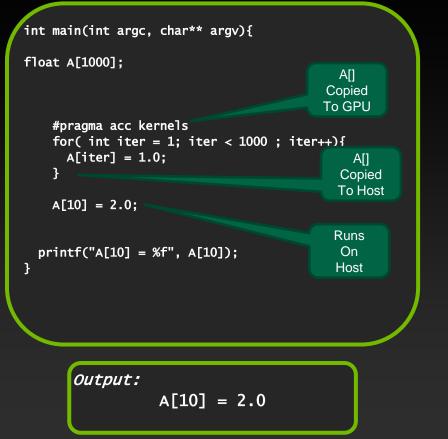
pgcc -acc -Minfo=accel loops.c
main:

- 6, Generating present\_or\_copyout(C[:])
  - \_Generating present\_or\_copy(B[:])
  - Generating present\_or\_copyout(A[:1000])
  - Generating NVIDIA code
- 7, Loop is parallelizable Accelerator kernel generated

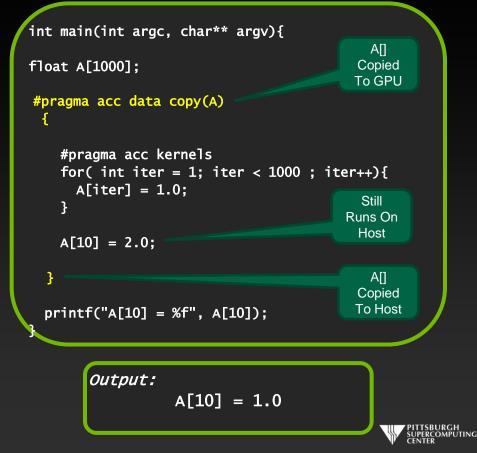


#### **Data Regions Have Real Consequences**

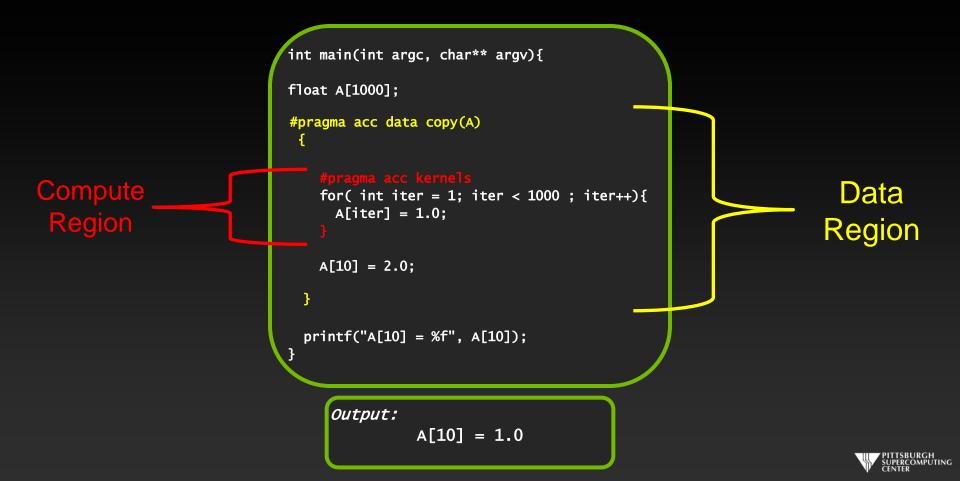
#### Simplest Kernel



#### With Global Data Region



#### Data Regions Are Different Than Compute Regions



#### **Data Movement Decisions**

- Much like loop data dependencies, sometime the compiler needs your human intelligence to make high-level decisions about data movement. Otherwise, it must remain conservative - sometimes at great cost.
- You must think about when data truly needs to migrate, and see if that is better than the default.
- Besides the scope based data clauses, there are OpenACC options to let us manage data movement more intensely or asynchronously. We could manage the above behavior with the update construct:

Fortran :
!\$acc update [host(), device(), ...]

C: #pragma acc update [host(), device(), ...]

Ex: #pragma acc update host(Temp\_array) //Gets host a current copy



# Exercise 2: Use acc data to minimize transfers

(about 40 minutes)

Q: What speedup can you get with data + kernels directives?

- Start with your Exercise 1 solution or grab laplace\_bad\_acc.c/f90 from the Solutions subdirectory. This is just the solution of the last exercise.
- Add data directives where it helps.
  - Think: when *should* I move data between host and GPU? Think how you would do it by hand, then determine which data clauses will implement that plan.
  - Hint: you may find it helpful to ignore the output at first and just concentrate on getting the solution to converge quickly (at 3372 steps). Then worry about *updating* the printout.



## Exercise 2 C Solution

#pragma acc data copy(Temperature\_last), create(Temperature)
while ( dt > MAX\_TEMP\_ERROR && iteration <= max\_iterations ) {</pre>

```
dt = 0.0; // reset largest temperature change
```

```
// copy grid to old grid for next iteration and find latest dt
#pragma acc kernels
for(i = 1; i <= ROWS; i++){
   for(j = 1; j <= COLUMNS; j++){
     dt = fmax( fabs(Temperature[i][j]-Temperature_last[i][j]), dt);
     Temperature_last[i][j] = Temperature[i][j];
   }
}</pre>
```

```
// periodically print test values
if((iteration % 100) == 0) {
    #pragma acc update host(Temperature)
    track_progress(iteration);
}
```



No data movement in this

block.



iteration++:

## **Exercise 2 Fortran Solution**

!\$acc data copy(temperature\_last), create(temperature)
do while ( dt > max\_temp\_error .and. iteration <= max\_iterations)</pre>

```
!$acc kernels
  do i=1.columns
     do i=1, rows
        temperature(i,j)=0.25*(temperature_last(i+1,j)+temperature_last(i-1,j)+ \&
                                temperature_last(i,j+1)+temperature_last(i,j-1))
     enddo
  enddo
   !$acc end kernels
  dt=0.0
   !copy grid to old grid for next iteration and find max change
  !$acc kernels
  do j=1.columns
     do i=1.rows
        dt = max( abs(temperature(i,j) - temperature_last(i,j)), dt )
        temperature_last(i,j) = temperature(i,j)
     enddo
  enddo
                                                           !$acc update host(temperature(columns-5:columns,rows-5:rows))
   !$acc end kernels
   !periodically print test values
                                                                                                         Except bring back a copy
  if( mod(iteration, 100).eq.0 ) then
      !$acc update host(temperature)
                                                                                                                      here
     call track_progress(temperature, iteration)
  endif
  iteration = iteration+1
enddo
```

!\$acc end data



Keep these on GPI

## Exercise 2: Performance

3372 steps to convergence

Execution	Time (s)	Speedup
CPU Serial (C)	36	
CPU 2 OpenMP threads	24	1.5x
CPU 4 OpenMP threads	15	2.4x
CPU 8 OpenMP threads	9.8	3.7x
CPU 16 OpenMP threads	5.0	7.2
OpenACC GPU	1.7	21x (3X vs. 16 CPU)

CPU: AMD 6276 Interlagos 8 Cores @ 2.3+GHz GPU: NVIDIA GK110 Kepler



## **Further speedups**

OpenACC gives us even more detailed control over parallelization
 Via gang, worker, and vector clauses

By understanding more about OpenACC execution model and GPU hardware organization, we can get higher speedups on this code

By understanding bottlenecks in the code via profiling, we can reorganize the code for higher performance

 But you have already gained most of any potential speedup, and you did it with a few lines of directives!



## **General Principles: Finding Parallelism In Code**

- Nested for/do loops are best for parallelization
  - Large loop counts are best
- Iterations of loops must be <u>independent</u> of each other
  - To help compiler: restrict keyword (C), independent clause
  - Use subscripted arrays, rather than pointer-indexed arrays (C)
- Data regions should avoid wasted transfers
  - If applicable, could use directives to explicitly control sizes
- Various other annoying things can interfere with accelerated regions
  - IO
  - Limitations on function calls and nested parallelism (relaxed much in 2.0)



## Is OpenACC Living Up To My Claims?

- High-level. No involvement of OpenCL, CUDA, etc.
- Single source. No forking off a separate GPU code. Compile the same program for accelerators or serial, non-GPU programmers can play along.
- Efficient. Experience show very favorable comparison to low-level implementations of same algorithms. kernels is magical!
- Performance portable. Supports GPU accelerators and co-processors from multiple vendors, current and future versions.
- Incremental. Developers can port and tune parts of their application as resources and profiling dictates. No wholesale rewrite required. Which can be <u>quick.</u>

